

§ Basics of Higher Cats III

§ 1 Recap on Straightening

why we need it:

most important tool in higher Cat. Theory; gives us

— lots of functors

— homotopy coherent algebras: sections + prop

$$R \mapsto R^{\text{Mod}}$$

$$\begin{array}{ccc} \text{Assoc} & \dashrightarrow & C^{\otimes} \\ & \searrow \cong & \downarrow \\ & & \text{Assoc} \end{array}$$

§ 2 Kan extensions

"extend" given functors universally

e.g.

$\mathcal{Q}\text{Cat}(-) : \text{PreStk} \longrightarrow \hat{\text{Cat}}$
needed heavily later on: convergent, (locally) of finite type prestacks
cocartesian prestacks

§ 3 Finality

Simplify computations for
(∞) limits, Kan extensions

S1 Straightening

Solva

Given $\Phi: E \rightarrow B$ we may consider
its fibres (for $b:B$)

$$\begin{array}{ccc} E_b & \longrightarrow & E \\ \downarrow & \lrcorner & \downarrow \\ \{b\} & \longrightarrow & B \end{array}$$

and ask: For which $\Phi: E \rightarrow B$
is

$$B \ni b \mapsto E_b \in \mathbf{Cat}$$

a functorial ... note that $\exists$ two
ways to be functorial: covariantly or
contravariantly

If $\Phi: E \rightarrow B$ is

- a Cartesian fibration $\Rightarrow$ $B \rightarrow \mathbf{Cat}$
 $b \mapsto E_b$

- a Cartesian fibration $\Rightarrow B^{\text{op}} \rightarrow \mathbf{Cat}$
 $b \mapsto E_b$

Straightening Unstraightening then says

$$(\mathbf{Cat}/B)_{\text{strict}} \begin{array}{c} \xrightarrow{\quad} \\ \sim \\ \xleftarrow{\quad} \end{array} \text{Fun}(B, \mathbf{Cat})$$

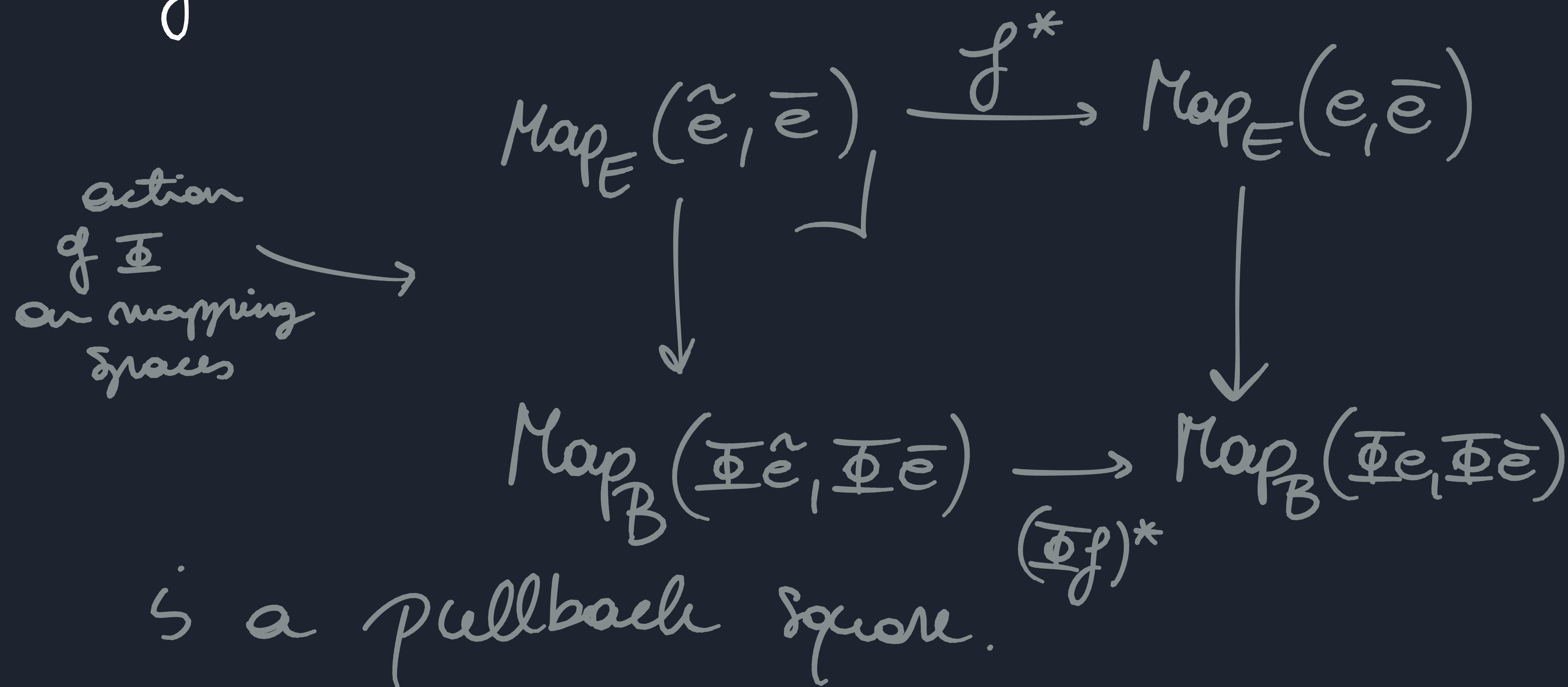
$$(\mathbf{Cat}/B)_{\text{flick}} \begin{array}{c} \xrightarrow{\quad} \\ \sim \\ \xleftarrow{\quad} \end{array} \text{Fun}(B^{\text{op}}, \mathbf{Cat})$$

Let us spell this out again:

def $\Phi: E \rightarrow B$

- A morphism $e \xrightarrow{f} \tilde{e}$ is called Φ -cartesian

if $\forall \bar{e} \in E$:



- A morphism $e \xrightarrow{f} \tilde{e}$ in E is called Φ -Cartesian if

$$\forall \bar{e} \in E$$

$$\text{Map}_E(\bar{e}, e) \xrightarrow{f_*} \text{Map}_E(\bar{e}, \tilde{e})$$



$$\text{Map}_B(\Phi \bar{e}, \Phi e)$$



$$\text{Map}_B(\Phi \bar{e}, \Phi \tilde{e})$$

is a pullback square.

Remark: $f: e \rightarrow \tilde{e}$ being Φ -CoCartesian means:

$$\begin{array}{ccc}
 e & \xrightarrow{f} & \tilde{e} \\
 \downarrow \scriptstyle \mathcal{H}h & \searrow \scriptstyle \tilde{\mathcal{I}}! & \downarrow \scriptstyle g_\alpha \\
 & \tilde{e} & \\
 & \downarrow & \\
 & e &
 \end{array}
 \quad \Phi g_\alpha \simeq \alpha$$

unique up to a contractible space of choices.

$$\begin{array}{ccc}
 & \downarrow & \\
 \Phi e & \xrightarrow{\Phi f} & \Phi \tilde{e} \\
 \downarrow \scriptstyle \Phi h & \searrow \scriptstyle \tilde{\mathcal{I}} & \downarrow \scriptstyle \mathcal{H}\alpha \\
 & \Phi e &
 \end{array}$$

Rule

Note that the definition of is not
a $\mathcal{D}$ -Cartesian morphism
Circular i.e

can define $f^*: \text{Map}_E(\tilde{e}, \bar{e}) \rightarrow \text{Map}_E(e, \bar{e})$
without straightening/unstraightening ...
Idea: Use "Axiom of ∞ -cats"

$$\text{Fun}([2], E) \xrightarrow{\sim} \text{Fun}([1], E) \times_E \text{Fun}([1], E)$$



def $\Phi: E \rightarrow B$ is called

- Φ -Cartesian fib. if for every $e \in E$ and every morphism $\beta: b \rightarrow \tilde{b}$ w/ $b \simeq \Phi e$,

there exists a Φ -Cartesian lift of β :

$$\text{lift}_0^\Phi(e, \beta): e \rightarrow \beta_! e$$

notation for the codomain

i.e. $\text{lift}_0^\Phi(e, \beta): e \rightarrow \beta_! e$ is Φ -Cart. and

$$\Phi(\text{lift}_0^\Phi(e, \beta)) \simeq \beta.$$

- Cartesian fibration if for all $\tilde{e} \in E$ and all $\beta: b \rightarrow \tilde{b}$ w/ $\tilde{b} \approx \Phi \tilde{e}$, there is a Φ -Cartesian lift

$$\text{lift}_1^\Phi(\tilde{e}, \beta) : \underbrace{\beta^* \tilde{e}}_{\text{notation for domain}} \rightarrow \beta$$

Remark Many equivalent (model independent) def of
(co)Cartesian fibrations ... see e.g. Riehl-Verity
Elements of ∞ -Cats ...

A conceptually very nice equivalent def
is : $\Phi: E \rightarrow B$ is a Cartesian fib

if $\vec{\text{ev}}_{\Phi}: E^{[1]} \rightarrow E \times_{\Phi} B$ has a
left adjoint section ... See also
Crossen
Cranshie etc.
Formalization
of Higher Cats

Recall For $B: 1\text{Cat}$, one defines

$$\begin{array}{ccc} 1\text{Cat}/B & \longrightarrow & \text{Fun}([1], 1\text{Cat}) \\ \downarrow & \lrcorner & \downarrow \\ * & \longrightarrow & 1\text{Cat} \end{array}$$

Objects : functors $E \rightarrow B$

morphisms :

$$\begin{array}{ccc} E & \xrightarrow{\psi} & E^{\sim} \\ & \searrow \quad \swarrow & \\ & B & \end{array}$$

def

- $\text{CoCart}/B \subset_{\text{full}} \text{Cat}/B$ full subcat.
on Cartesian fibrations on B .

- $\text{Cat}/B \subset_{\text{full}} \text{Cat}/B$ analogously...

- $(\text{CoCart}/B)_{\text{strict 1-full}} \subset \text{CoCart}/B$

1-full subcategory on those morphisms

$$\begin{array}{ccc} E & \xrightarrow{\psi} & \tilde{E} \\ \Phi \searrow & \tilde{\sim} & \swarrow \tilde{\Phi} \\ & B & \end{array}$$

such that ψ
sends Φ -Cartesian
morphisms to
 $\tilde{\Phi}$ -Cartesian
morphisms.

- Analogously: $(\text{Cart}/B)_{\text{strict}} \subset_{1\text{-full}} 1\text{Cart}/B$
- $0\text{-CoCart}/B \subset_{\text{full}} \text{CoCart}/B$ on these colimitation fibrations $E \xrightarrow{\Phi} B$ such that E_b is a space for every $b \in B$.

Note:

$0\text{-CoCart}/B \cap (\text{CoCart}/B)_{\text{strict}}$
 $\downarrow \rho$ is an equivalence
 $0\text{CoCart}/B$

-
- Analogously: $0\text{Cart}/B \subset_{\text{full}} \text{Cart}/B$

Then

There are equivalences

$$\begin{array}{ccc} \text{Fun}(B, \mathbf{1Cat}) & \begin{array}{c} \xrightarrow{\text{Un}} \\ \xleftarrow[\cong]{} \end{array} & (\mathbf{CoCat}/B)_{\text{strict}} \\ \downarrow \text{IU} & & \downarrow \text{IU} \\ \text{Fun}(B, \mathbf{Spec}) & \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow[\sim]{} \end{array} & \mathbf{O-CoCat}/B \end{array}$$

Contravariantly natural in B w.r.t/
precomposition on the left and pullback
on the right.

Passing to opposite Cats:

$$\text{Fun}(B^{\text{op}}, 1_{\text{Cat}}) \begin{array}{c} \xrightarrow{Un} \\ \xleftarrow{\sim} \\ \text{st} \end{array} (\text{Cat}/B)_{\text{strict}} \quad \begin{array}{c} \\ \\ IU \end{array}$$

$$\text{Fun}(B^{\text{op}}, \text{Spec}) \begin{array}{c} \xrightarrow{\sim} \\ \xleftarrow{\sim} \end{array} 0\text{-Cat}/B \quad \begin{array}{c} \\ \\ IU \end{array}$$

More explicitly:

Given
obtain

$$\Phi: E \rightarrow B$$

coCartesian fibration, we

$$\begin{array}{ccc} B & \xrightarrow{\mathcal{S}\Phi} & \mathbf{Cat} \\ b & \mapsto & E_b \\ \downarrow \mathcal{P} & \mapsto & \downarrow \mathcal{P}_! \\ \tilde{b} & \mapsto & E_{\tilde{b}} \end{array}$$

action on morphism spaces is induced by
 Φ -coCartesian lifts.

Conversely: A functor
 $F: B \rightarrow \mathbf{Cat}$

gives rise to

$$\begin{array}{ccc} b & \mapsto & Fb \\ \beta \downarrow & \mapsto & \downarrow F_\beta \\ \tilde{b} & \mapsto & F\tilde{b} \end{array}$$

a CoCartesian fib.

$$\mathrm{Un}_F: B^{\circ} \longrightarrow B$$

where B° is informally described by:

- objects: $(b \in B, x: Fb) \iff \begin{array}{c} * \\ \downarrow^x \\ Fb \end{array}$

- morphisms:

$$\begin{array}{ccc} \begin{array}{c} b \\ \beta \downarrow \\ \tilde{b} \end{array} & \begin{array}{ccc} x & * & \tilde{x} \\ \swarrow & \searrow & \downarrow \\ Fb & \xrightarrow{F_\beta} & F\tilde{b} \end{array} & \iff & \begin{array}{ccc} F_\beta(x) & \xrightarrow{\sim} & \tilde{x} \\ \uparrow (\mathrm{Un}_F - \text{CoCart}) & & \downarrow \text{iso} \end{array} \end{array}$$

The functor

$$\begin{array}{ccc} \mathrm{Un}_F: B^{\circ} & \longrightarrow & B \\ (b, x) & \longmapsto & b \end{array}$$

does the obvious thing.

Prop There exists a universal Cartesian fib.

1Cat.



1Cat

such that

$$\begin{array}{ccc} B^{\mathcal{F}} & \xrightarrow{\quad} & 1\text{Cat.} \\ \text{Un}_F \downarrow \lrcorner & & \downarrow \pi^{\text{univ}} \\ B & \xrightarrow{\quad F \quad} & 1\text{Cat} \end{array}$$

"Any Cartesian fib. is obtained as a pullback from π^{univ} "

Examples:

- C cat, $c \in C$

- $C_{c/} \rightarrow C$ coCartesian fib

$\downarrow \text{st}$

$C(c, -) : C \rightarrow \text{Spc}$

$C_{/c} \rightarrow C$ cartesian fib

$\downarrow \text{st}$

$C(-, c) : C^{\text{op}} \rightarrow \text{Spc}$

- monoidal cats will be certain coCartesian fib

$C^{\otimes} \rightarrow \Delta^{\text{op}}$

- C monoidal category, C^+ a left C -module category.

One defines a Cartesian fibration

$\text{Alg} + \text{Mod}(C, C^+) \rightarrow \text{Alg}(C)$

which straightens to

$\text{Alg}(C)^{\text{op}} \xrightarrow{\quad} \widehat{\text{Cat}}$

A

$\hookrightarrow$

$A \text{Mod}(C^+)$

Example (1.5 Yoneda) ← will not can this in my talk
 in the code

Note that both

$$\begin{array}{ccc} \text{Fun}([1], C) & & C \times C \\ \text{ev}_0 \downarrow & \text{and} & \downarrow \text{proj}_1 \\ C & & C \end{array}$$

are cartesian fibrations.

An ev_0 -cartesian morphism in $\text{Fun}([1], C)$ is a morphism $d: f \rightarrow f'$ s.t.

$$\begin{array}{ccc} c_0 & \xrightarrow{d_0} & c'_0 \\ f \downarrow & \simeq & \downarrow f' \\ c_1 & \xrightarrow{d_1} & c'_1 \end{array}$$

$\text{ev}_1(d) = d_1$ is an isomorphism.

Analogously: proj_1 -cartesian morphism is a morphism $(d_0, d_1): (c_0, c_1) \rightarrow (c'_0, c'_1)$ s.t. $\text{proj}_2(d_0, d_1) = d_1$ is an iso.

In particular, the commutative triangle

$$\begin{array}{ccc} \text{Fun}([1], C) & \xrightarrow{(\text{ev}_0, \text{ev}_1)} & C \times C \\ & \searrow \text{ev}_0 & \swarrow \text{proj}_1 \\ & C & \end{array}$$

is a morphism in (Cart/C) which $\simeq \text{Fun}(C^{\text{op}}, 1\text{Cat})$

$(\text{ev}_0, \text{ev}_1)$ clearly sends

ev_0 -cartesian morphisms to proj_1 -cartesian morphisms

Straightening this morphism yields a natural transformation

$$\textcircled{*} (c \mapsto C_c) \implies \subseteq$$

components are

$$C_c \xrightarrow{\text{ev}_c} C$$

by def of straightening

constant diagram
functor

$$C^{\text{op}} \rightarrow \mathbf{Cat}$$

$$x \mapsto C$$

will not
speak about
this

Note that:

$$\text{Fun}(C^{\text{op}}, \mathbf{Cat}) / \cong \xrightarrow{\sim} \text{Fun}(C^{\text{op}}, \mathbf{Cat}/C)$$

and hence

$\textcircled{*}$ gives us a functor

$$\textcircled{**} C^{\text{op}} \rightarrow \mathbf{Cat}/C$$

$$c \mapsto (C_c \rightarrow C)$$

O -CoCartesian fib
w/ fiber $\text{Map}_C(c, x) \in \mathbf{Spc}$
at x

and therefore $\textcircled{**}$ factors as

$$C^{\text{op}} \rightarrow \mathbf{Cat}/C$$

$$\downarrow \cong$$

$$O\text{-CoCat}/C$$

$$\cong \text{Fun}(C, \mathbf{Spc})$$

So by currying we
get the Hom functor $C^{\text{op}} \times C \xrightarrow{\text{Map}_C} \mathbf{Spc}$.

1.7 Adjoint functors

$F: C_0 \rightarrow C_1$ functor

$$\iff [1] \xrightarrow{F} 1_{\text{Cat}}$$

unstraighten

$$\iff \tilde{C} \xrightarrow{\tilde{F}} [1] \quad \text{Cartesian fibration}$$

def $F: C_0 \rightarrow C_1$ admits a right adjoint
if $\tilde{C} \xrightarrow{\tilde{F}} [1]$ is also a Cartesian
fibration (i.e. if is bicartesian).

Assume $F: C_0 \rightarrow C_1$ admits a RA. Then we may apply cartesian unbraiding to $\hat{C} \xrightarrow{\hat{F}} [1]$ to obtain

$$\begin{array}{ccc}
 [1]^{\text{op}} & \longrightarrow & 1_{C_0} \\
 \uparrow & \xrightarrow{\quad} & \uparrow \\
 0 & \xrightarrow{\quad} & C_0 \\
 \uparrow & \xrightarrow{\quad} & \uparrow F^R \\
 1 & \xrightarrow{\quad} & C_1
 \end{array}$$

hence we get a functor $C_1 \xrightarrow{F^R} C_0$.

Exercise: Show that for all $x \in C_0$ $y \in C_1$ one has

$$\text{Map}_{C_1}(Fx, y) \cong \text{Map}_{C_0}(x, F^R y)$$

1.7.3 Partial Adjoints

$\tilde{C} \xrightarrow{\tilde{F}} [1]$ as before without assuming $\tilde{F}$ is a Cartesian fibration.

Let $C'_1 \subset C_1$ full subcat. on those

$e_1 \in C_1$ such that there exists an $\tilde{F}$ -Cartesian morphism

$c_0 \rightarrow c_1$ in $\tilde{C}$ covering $0 \rightarrow 1$ in $[1]$.

Let $\tilde{C}' \subset \tilde{C}$ full subcat so that

$$\tilde{C}'_0 = C_0$$

$$\tilde{C}'_1 = C'_1$$

Then :

$\tilde{C}' \rightarrow [1]$ is a Cartesian fibration

$\Rightarrow$ partially defined right adjoint

$$C'_1 \xrightarrow{F'^R} C_0$$

$$\text{Map}_{C'_1}(F C_0, c_1) \simeq \text{Map}_{C_0}(c_0, F'^R c_1)$$

for all $c_0 \in C_0$
 $c_1 \in C'_1$

2. Basic Operations on ∞ -Cats

Q1 How to produce new ∞ -Cats from old ones? (e.g. derived Cats)

Q2 How to construct new functors from old ones?

A1 (co)limits e.g. localizations
→ derived ∞ -Cat

realised as a pushout

A2 Kan extensions e.g. construction of $2\text{Cat}(-)$

2.1.1

$$F: D \rightarrow C$$

For any category E we have

$$\text{Fun}(C, E) \xrightarrow{F^*} \text{Fun}(D, E)$$

precomposition functor

2.1.2

def

partially defined left (resp. right)

adjoint of F^* is called

left (resp. right) Kan extension

along F and denoted by

LKE_F (resp. RKE_F)

Example we will see many examples of
Kan extensions in future sessions;

Easy examples are given by

(co) limits:

If $C = *$, $F = (D \xrightarrow{!} *)$

Then left and right Kan extension
give

$$\operatorname{Colim}_D : \operatorname{Fun}(D, E) \longrightarrow E$$

$$\operatorname{lim}_D : \operatorname{Fun}(D, E) \longrightarrow E$$

(partially defined)

2.1.3 Formulas for Kan extensions

Let

$$\begin{array}{ccc} D & \xrightarrow{\Phi} & E \\ F \downarrow & & \\ C & & \end{array}$$

Consider for $c \in C$

$$\begin{array}{ccc} F \downarrow c \simeq D \times_C C/c & \longrightarrow & C/c \\ \pi_D \downarrow & \lrcorner & \downarrow \\ D & \xrightarrow{F} & C \end{array}$$

objects: $Fd \xrightarrow{\alpha} c$

morphisms:

$$\begin{array}{ccc} 1 & d & \\ & \downarrow & \\ & \tilde{d} & \end{array} \quad \begin{array}{ccc} Fd & \xrightarrow{F_1} & F\tilde{d} \\ \alpha \searrow & \simeq & \swarrow \tilde{\alpha} \\ & c & \end{array}$$

and suppose that for every $c \in C$

$$\text{notation } \left\{ \begin{array}{c} \text{Colim} \\ D \times_C C/c \end{array} \right. \Phi := \text{Colim}_{D \times_C C/c} (\Phi \circ \pi_D) \text{ exists.}$$

Then: $LK_{E_F} \Phi : C \rightarrow E$ exists and
 an object is given by

$$c \mapsto \operatorname{colim}_{D \times_C C/c} \Phi \simeq LK_{E_F} \Phi(c)$$

Similarly:

$$RK_{E_F} \Phi(c) \simeq \lim_{D \times_C C/c} \Phi$$

Exercise:

$$\begin{aligned} \operatorname{colim}_D \Phi &\simeq \operatorname{colim}_C LK_{E_F}(\Phi) \\ \lim_D \Phi &\simeq \lim_C RK_{E_F}(\Phi) \end{aligned}$$

2.1.5 There is an adjunction

$$\mathcal{S}pc \begin{matrix} \xleftarrow{|-|} \\ \xrightarrow{\perp} \\ \xleftarrow{\perp} \\ \xrightarrow{(-)^{\approx}} \end{matrix} \mathcal{I}Cat$$

maximal
subspace

Note that

$$|C| \approx \operatorname{colim}_C *$$

constant
diagram
functor
 $C \rightarrow \mathcal{S}pc$
 $c \mapsto *$

For any
 $S \in \mathcal{S}pc$ we
have

$$|S| \approx S.$$

In fact, this may be proven
by straightening-unstraightening:

$$\operatorname{Map}_{\mathcal{S}pc}(\operatorname{colim}_C *, S)$$

$$\approx \lim_C \operatorname{Map}_{\mathcal{S}pc}(*, S)$$

$$\approx \lim_C \underline{S} \approx \Gamma(C \times S \rightarrow C)$$

$$\approx \operatorname{Fun}(C, S)$$

$$\approx \operatorname{Map}_{\mathcal{S}pc}(C, S)$$

as wanted

will not
come up in
lecture

def $C: \mathbf{Cat}$ is called contractible if

$$|C| \simeq *$$

2.1.6 If $\underline{\Phi} : C \rightarrow \mathbf{Spc}$, then

$$\operatorname{colim}_C \underline{\Phi} \simeq |\tilde{C}_{\underline{\Phi}}| \quad \text{where}$$

$\tilde{C}_{\underline{\Phi}} \rightarrow C$ is the coCartesian fib. assoc.
to $\underline{\Phi}$.

Lemma

$$C \in 1\text{Cat}, \quad PC := \text{Fun}(C^{\text{op}}, \text{Spc})$$

Then restriction and left Kan extension
along the Yoneda embedding

$$C \xrightarrow{Y} PC$$

induce for all cocomplete D an
equivalence

$$\text{Fun}^{\text{colim}}(PC, D) \xrightleftharpoons[\text{LKE}_Y]{Y^*} \text{Fun}(C, D)$$

exhibiting PC as the free cocompletion of C .

In particular: $\text{Fun}^{\text{colim}}(\text{Spc}, D) \simeq D$.

§3 Cofinality

2.2.1

def

$F: D \rightarrow C$ is called cofinal if for any $c: C$, the category

$c \downarrow F \simeq D \times_C c$ is contractible i.e.

$$|D \times_C c| \simeq *$$

(and coinital if $F^{op}: D^{op} \rightarrow C^{op}$ is cofinal)
i.e. $|F \downarrow c| \simeq *$

2.2.2

Lemma TFAE:

(i) $F: D \rightarrow C$ is cofinal.

(ii) For any $\Phi: C \rightarrow E$, the natural map

$$\operatorname{colim}_D (\Phi \circ F) \xrightarrow{\sim} \operatorname{colim}_C \Phi$$

is an isomorphism.

(iii) For any $\Phi: C^{\text{op}} \rightarrow E$, the natural map

$$\lim_{C^{\text{op}}} \Phi \xrightarrow{\sim} \lim_{D^{\text{op}}} (\Phi \circ F^{\text{op}})$$

is an iso.

for the book there is also: (ii'), (ii''), (iii'), (iv)

2.2.3

Example

- Let C be a category w/ a terminal object $t \in C$. Then the functor

$$* \xrightarrow{t} C \text{ is cofinal i.e.}$$

for $\Phi: C \rightarrow E$ we have

$$\operatorname{colim}_C \Phi \approx \operatorname{colim}_* \Phi(t) \approx \Phi(t)$$

- More generally: Any functor

$$F: D \rightarrow C \text{ that admits a}$$

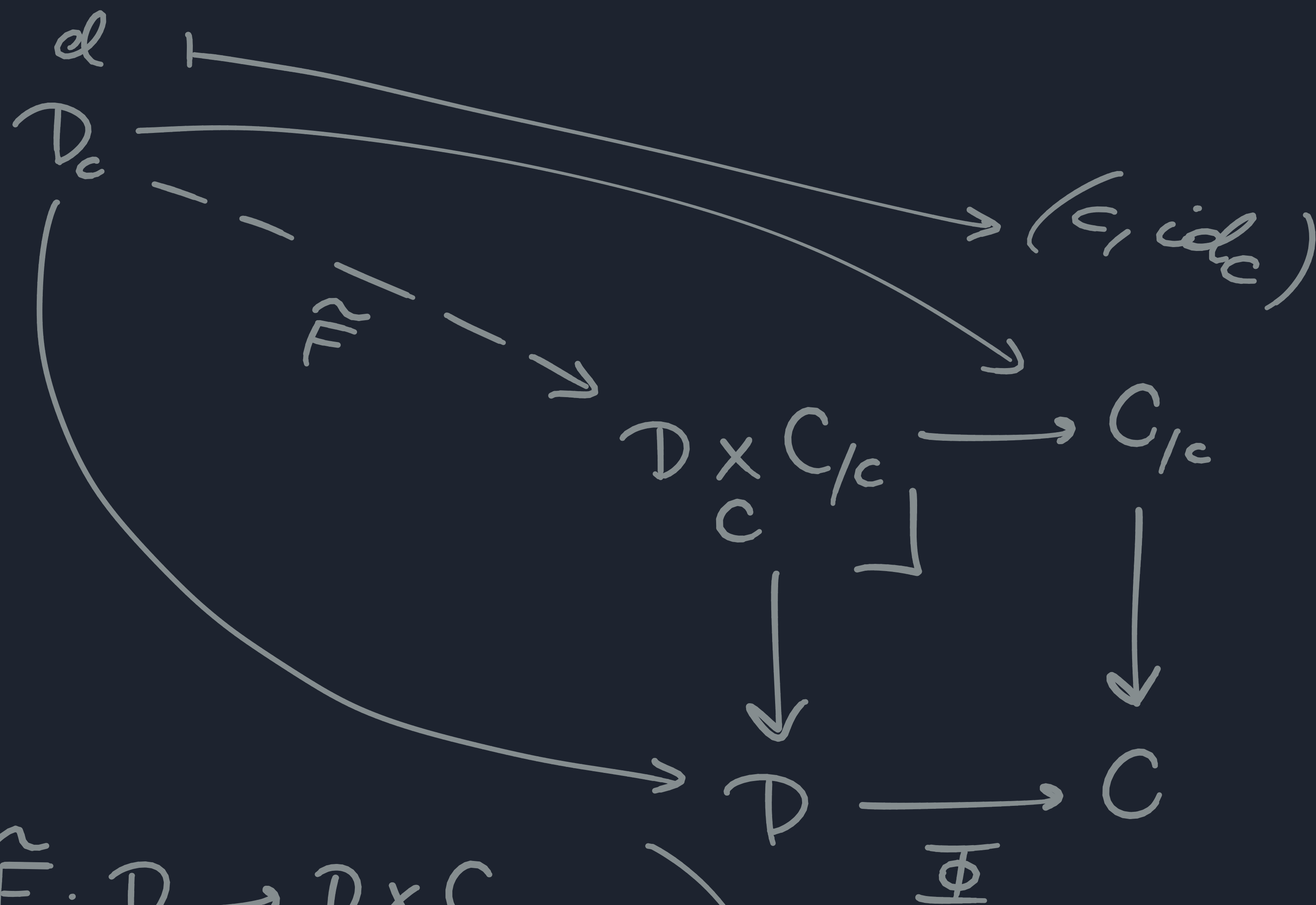
left adjoint (i.e. F is a right adjoint) is

cofinal.

Indeed: $D \times_C C_c \approx c \downarrow F$ has an initial object for all c
and hence $|D \times_C C_c| \approx *$.

Example/Lemma

$D \xrightarrow{F} C$ is Cartesian fibration
Then the canonical functor



(i.e. $\tilde{F}: D_c \rightarrow D \times_C C/c$
 $d \mapsto (d, id_c)$)

is cofinal.

See Kerdan Prop 7.3.4.1

Hence we get:

$$LKE_F \Phi(c) \simeq \operatorname{Colim}_{D \times C/c} \Phi \simeq \operatorname{Colim}_{D_c} \Phi$$

Similarly, if $D \xrightarrow{F} C$ is a Cartesian fib.

$$RKE_F \Phi(c) \simeq \operatorname{lim}_{D_c} \Phi$$

2.3 The book also says something on
contraction functors ...

we will skip this for now